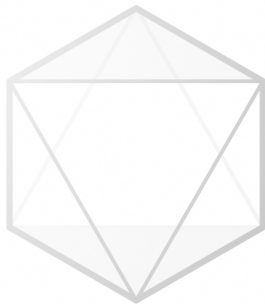
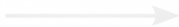


THE STRUCTURE OF 2-LEVEL POLYTOPES

Martin Winter

TU Berlin

22. January, 2025

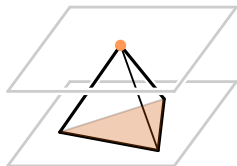
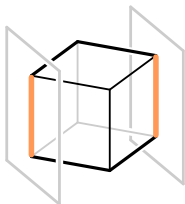
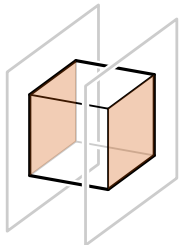


2-LEVEL POLYTOPES

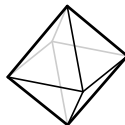
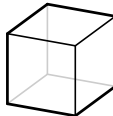
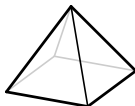
$$P = \text{conv}\{p_1, \dots, p_n\} \subset \mathbb{R}^d, \quad d \geq 0$$

Definition.

- ▶ Two faces $F_1, F_2 \subseteq P$ are **antipodal** if they are contained in parallel hyperplanes (i.e. there are parallel hyperplanes $H_1, H_2 \subseteq \mathbb{R}^d$ with $F_i = P \cap H_i$)
- ▶ A polytope P is **2-level** if each antipodal face pair that contains a facet also contains all vertices.



EXAMPLES



dim	0	1	2	3	4	5	6	7	8
2-level	1	1	2	5	19	106	1150	27291	1378453

EXAMPLES

Many 2-level polytopes are constructed from combinatorial objects:

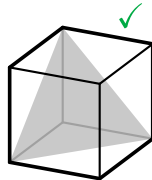
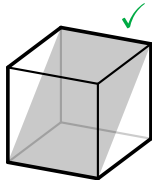
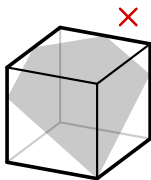
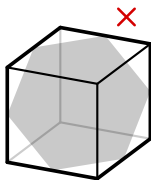
- ▶ *Hanner polytopes* (in relation to cographs)
- ▶ order polytopes of posets
- ▶ stable set polytopes of perfect graphs
+ their twisted prisms (= *Hansen polytopes*)
- ▶ spanning tree polytopes of series-parallel graphs
- ▶ *Birkhoff polytopes* (from double stochastic matrices)
- ▶ certain matroid base polytopes

PROPERTIES

- ▶ all faces are 2-level
- ▶ closed under products and joins
- ▶ $\#\text{vertices} \cdot \#\text{facets} \leq d2^{d+1}$
- ▶ are 01-polytopes (if P is d -dimensional then $P \subseteq [0, 1]^d$)

Theorem.

2-level polytopes are precisely the polytopes that can be written as the intersection of a cube with an affine subspace that is spanned by vertices of the cube.



BINARY SCALAR PRODUCTS

Based on empirical evidence it was conjectured that

$$f_0(P) \cdot f_{d-1}(P) \leq d2^{d+1}.$$

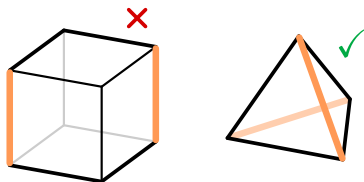
This was eventually proven using **binary scalar product pairs**: $A, B \subset \mathbb{R}^d$

$$\langle a, b \rangle \in \{0, 1\}, \quad \text{for all } a \in A \text{ and } b \in B.$$

Theorem. (KUPAVSKII, WELTGE; 2020)

If $A, B \subseteq \mathbb{R}^d$ is a binary scalar product pair, then $|A| \cdot |B| \leq (d + 1)2^d$.

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$\mathcal{A}_s(P)$... set of spanning antipodal face pairs

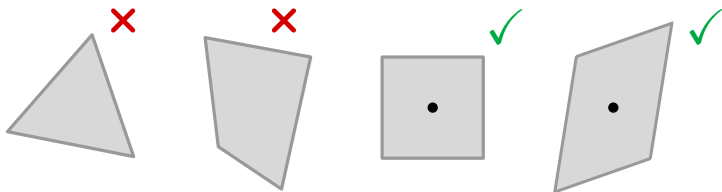
Corollary.

For a 2-level polytope holds $f_0(P) \cdot |\mathcal{A}_s(P)| \leq (d + 1)2^d$.

CONJECTURES FOR CENTRALLY
SYMMETRIC POLYTOPES

CENTRALLY SYMMETRIC POLYTOPES

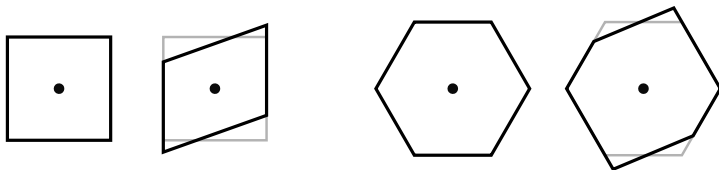
centrally symmetric $:\Leftrightarrow P = -P$



dim	0	1	2	3	4	5	6	7	8
2-level	1	1	2	5	19	106	1150	27291	1378453
cs 2-level	1	1	1	2	4	13	45	238	1790

PROPERTIES

- ▶ each facet contains exactly half of the vertices
 - ▶ each vertex lies in exactly half of the facets
- } defining 2-level among cs
- ⇒ closed under polar duality
 - ▶ precisely the section of a cube with linear space spanned by vertices
 - ▶ linearly unique (all cs 01-polytopes are linearly unique)
 - ⇒ linearly invariant geometric properties are combinatorial properties



KALAI'S 3^d CONJECTURE

$$s(P) := f_{-1} + f_0 + f_1 + \cdots + f_{d-1} + f_d = \#\text{non-empty faces}$$

Conjecture. (3^d conjecture, KALAI, 1989)

For every centrally symmetric d -polytope $P \subset \mathbb{R}^d$ holds

$$s(P) \geq s(d\text{-cube}) = 3^d.$$

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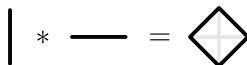
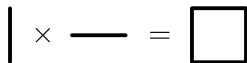
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But: cube is not the only minimizer! \rightarrow **Hanner polytopes**

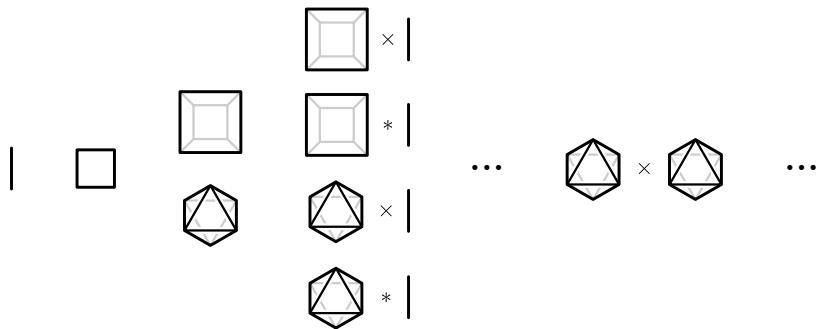
HANNER POLYTOPES

Hanner polytopes are defined recursively:

- (i) start from a line segment.
- (ii) recursively apply Cartesian products (\times) and sums ($*$)



HANNER POLYTOPES



#Hanner polytopes for $d \geq 1 = 1, 1, 2, 4, 8, 18, 40, 94, 224, 548, \dots$

KALAI'S 3^d CONJECTURE

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$$s(P) \geq s(d\text{-cube}) = 3^d.$$

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What is known ... ?

- ▶ dimension $d \leq 3$ ✓ easy
- ▶ dimension $d = 4$ ✓ not so easy (2007)
- ▶ simple/simplicial polytopes ✓ needs a lot of algebra
- ▶ without requiring central symmetry ✓ easy $\rightarrow s(d\text{-simplex}) = 2^d - 1$

MAHLER'S CONJECTURE

Mahler volume ... $M(P) := \text{vol}(P) \cdot \text{vol}(P^\circ)$

Conjecture. (3^d conjecture, MAHLER, 1939)

For every centrally symmetric d -polytope $P \subset \mathbb{R}^d$ holds

$$\text{measures "roundness"} \longrightarrow M(P) \geq M(d\text{-cube}) = \frac{4^d}{d!}.$$

But: cube is not the only minimizer! \rightarrow **Hanner polytopes**

What is known ... ?

- ▶ dimension $d \leq 3$ ✓ not so easy ($d = 2$: 1939, $d = 3$: 2020)
- ▶ dimension $d = 4$? out of reach
- ▶ Hanner polytopes are local minimizers ✓ (2014)
- ▶ without requiring central symmetry ? open $\rightarrow M(d\text{-simplex}) = \frac{(d+1)^{d+1}}{(d!)^2}$

KALAI'S FLAG CONJECTURE

$$S(P) := \# \text{flags of } P$$

Conjecture. (flag conjecture, KALAI, 1989)

For every centrally symmetric d -polytope $P \subset \mathbb{R}^d$ holds

$$S(P) \geq S(d\text{-cube}) = d! 2^d.$$

But: cube is not the only minimizer! \rightarrow **Hanner polytopes**

KALAI'S 3^d CONJECTURE

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- ▶ verified for $d \leq 7$ by enumeration

KALAI'S 3^d CONJECTURE

- ▶ verified for $d \leq 7$ by enumeration
- ▶ 2-level polytopes are “very small polytopes”

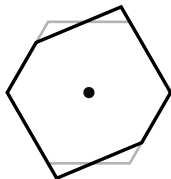
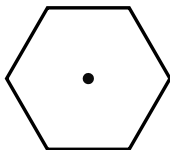
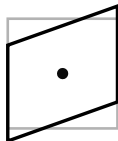
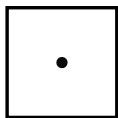
LINEARLY UNIQUE POLYTOPES

Let $\text{REAL}(P)$ be the space of cs realization of P module linear transformations.

Definition.

A centrally symmetric polytope is

- ▶ **linearly unique** if $\text{REAL}(P)$ consists of a single point.
- ▶ **linearly discrete** if $\text{REAL}(P)$ consists of finitely many points.
- ▶ **linearly compact** if $\text{REAL}(P)$ is compact.



LINEARLY COMPACT POLYTOPES

Lemma.

If P is a cs minimizer of face number, then P is linearly compact.

Proof sketch.

- ▶ if P is not linearly compact, then there is a convergent sequence P_1, P_2, P_3, \dots of realizations of P with $\lim P_n$ not being a realization of P .
- ▶ observe that in the limit, there cannot be new faces, but faces must have vanished.
 $\implies P$ cannot have been a minimizer. □

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Conjecture.

The only polytopes with compact realization spaces are linearly discrete.

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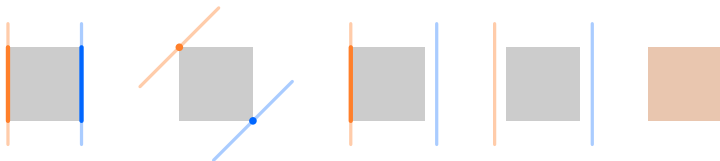
Conjecture.

The only polytopes with compact realization spaces are linearly discrete.

- ▶ true for matroids and oriented matroids.
- ▶ polytope realization spaces are unions of oriented matroid realization spaces.

KALAI FOR NON-CS POLYTOPES

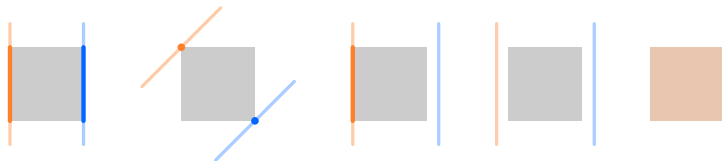
- ▶ By convention, say that (P, \emptyset) is an antipodal face pair as well.



- ▶ $\mathcal{A}(d)$... set of ordered antipodal face pairs
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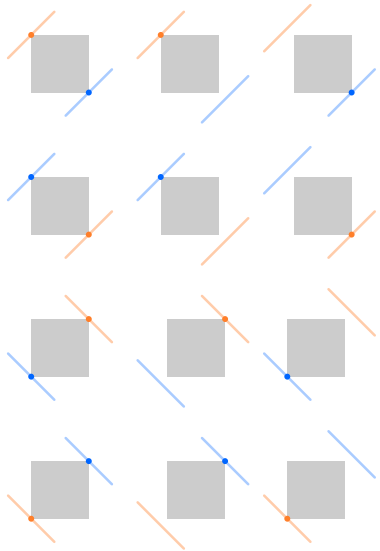
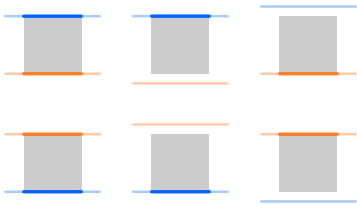
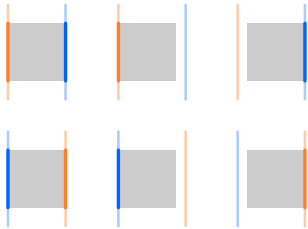
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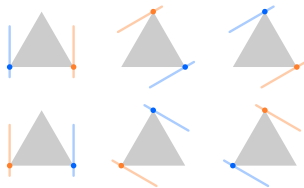
Conjecture.

If P is a (2-level) polytope, then

$$|\mathcal{A}(P)| \geq 3^{d+1}$$

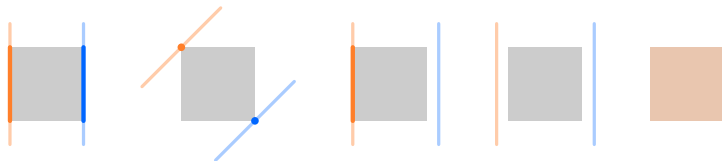
$$|\mathcal{A}^*(P)| \geq \frac{3^{d+1} + 1}{2} = 3 \cdot \frac{3^d - 1}{2} + 2$$





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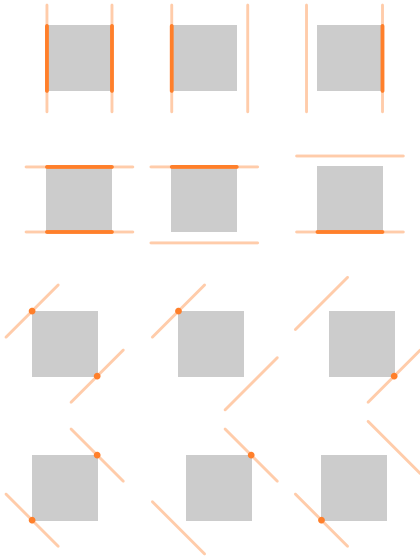
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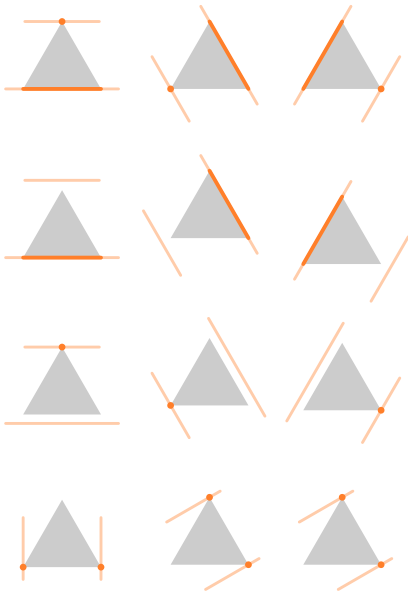
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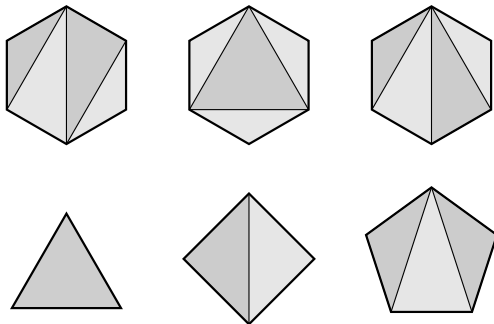




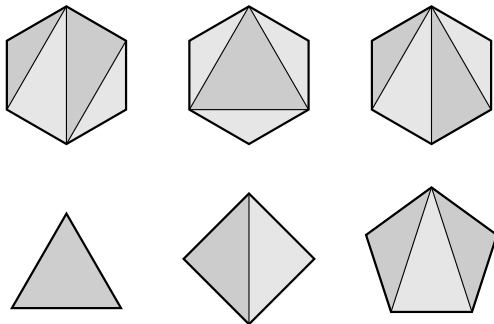
MAHLER'S CONJECTURE

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VOLUME AND TRIANGULATION



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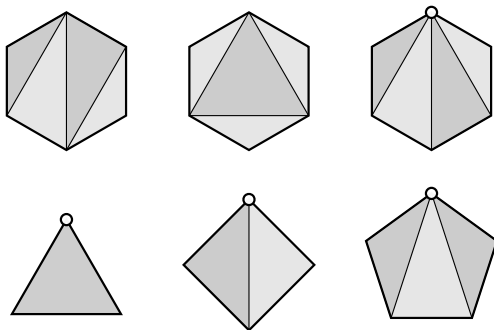


Theorem.

In a 2-level polytopes ...

- ▶ *each simplex in a pulling triangulation has the same volume.*
(lattice volume 1)

VOLUME AND TRIANGULATION

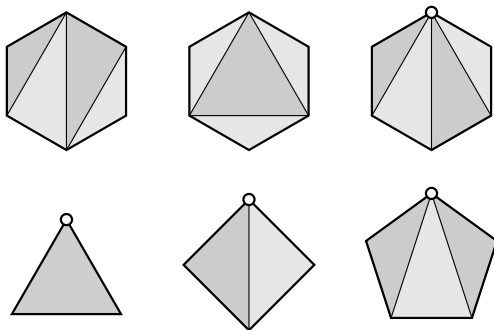


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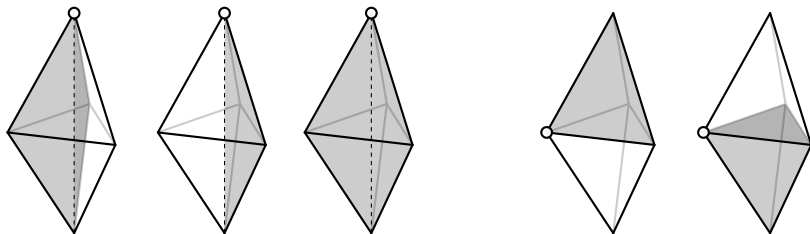


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MAHLER VOLUME AND PULLING TRIANGULATIONS

$f_d^*(P)$... # simplices in pulling triangulation of P

The Mahler conjecture is equivalent to the following:

Conjecture.

For a centrally-symmetric 2-level polytope $P \subset \mathbb{R}^d$ holds

$$f_d^*(P) \cdot f_d^*(P^\circ) \geq d!2^{d-1}$$

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↑
flag conjecture

PULLING TRIANGULATIONS

Theorem.

If P is a 2-level polytope, then

$$f_d^*(P) = \sum_{F_0 \preceq \dots \preceq F_d} \left(1 - \frac{f_0(F_0)}{f_0(F_1)}\right) \cdots \left(1 - \frac{f_0(F_{d-1})}{f_0(F_d)}\right).$$

Proof.

- ▶ For a vertex v and facet F_{d-1} , let $[v \notin F_{d-1}]$ denote the indicator function.
- ▶ We have

$$f_d^*(P) = \sum_{F_{d-1}} [v \notin F_{d-1}] f_{d-1}^*(F_{d-1}).$$

- ▶ Take expectation value *w.r.t.* a uniform random choice of v :

$$f_d^*(P) = \sum_{F_{d-1}} \left(1 - \frac{f_0(F_{d-1})}{f_0(P)}\right) f_{d-1}^*(F_{d-1}) = \sum_{F_0 \preceq \dots \preceq F_d} \left(1 - \frac{f_0(F_0)}{f_0(F_1)}\right) \cdots \left(1 - \frac{f_0(F_{d-1})}{f_0(F_d)}\right)$$

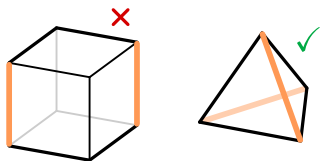
□

STRONGLY 2-LEVEL POLYTOPES

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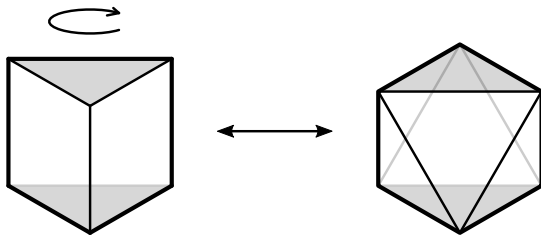
Definition.

- ▶ An antipodal face pair is **spanning** if its affine span is \mathbb{R}^d .
- ▶ A polytope P is **strongly 2-level** if each spanning antipodal face pair contains all vertices.

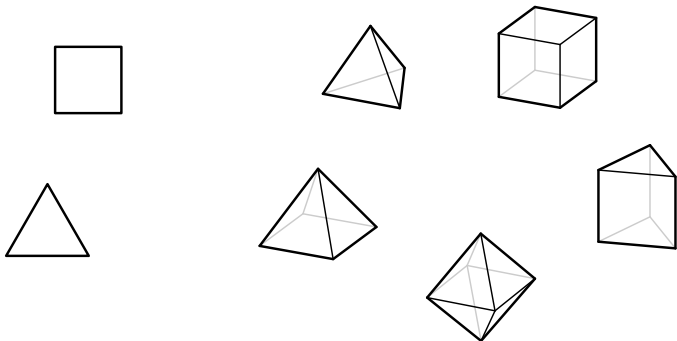


- ▶ $d \leq 4$: every 2-level polytope is strongly 2-level.
- ▶ $d = 5$: there is a unique 2-level polytope that is not strongly 2-level (the (6,2)-hypersimplex)
- ▶ $d \leq 5$: faces of strongly 2-level polytopes are strongly 2-level.
- ▶ $d = 6$: there is a unique strongly 2-level polytope with a facet that is not. (the (7,3)-hypersimplex with a (6,2)-hypersimplex as a facet)

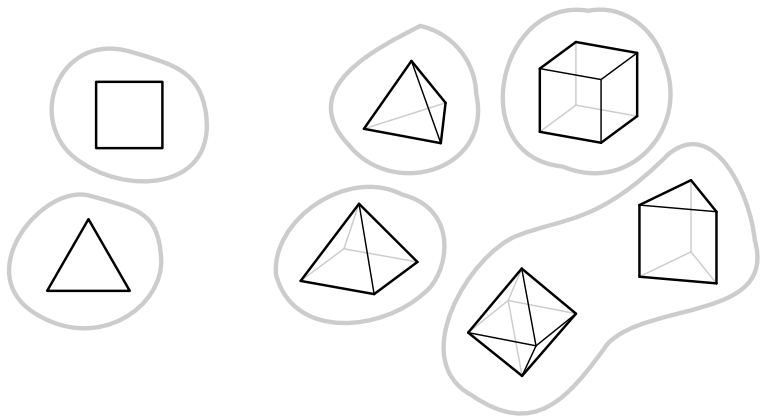
TWISTING



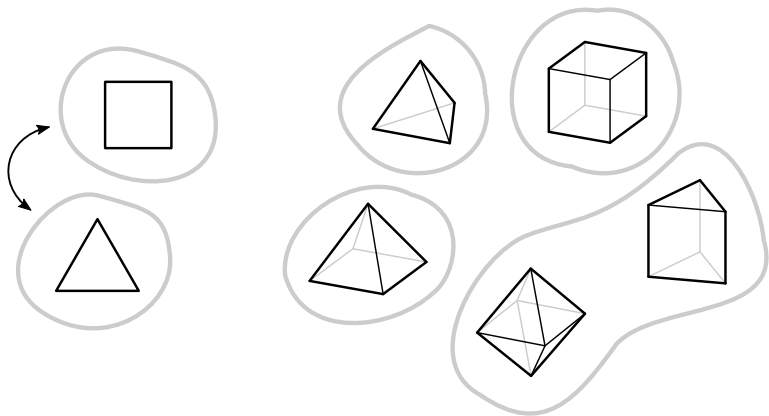
TWISTING CLASSES



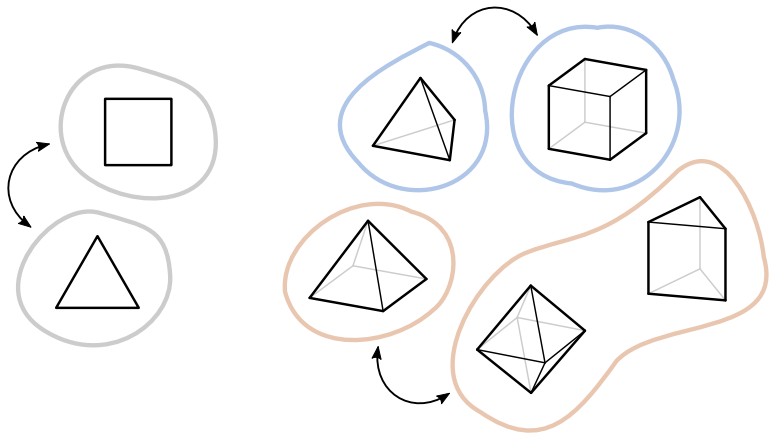
TWISTING CLASSES



TWISTING CLASSES & TWISTING DUALITY



TWISTING CLASSES & TWISTING DUALITY



TWISTING DUALITY

All members of a twisting class share

- ▶ # vertices
- ▶ # vertex orbits
- ▶ # antipodal face pairs
- ▶ # spanning antipodal face pairs

Within a twisting class one can get from anywhere to anywhere by a single twist.

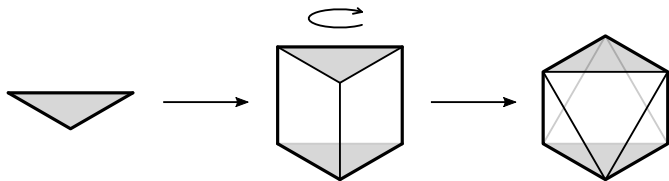
Twisting duality swaps:

- ▶ # vertex orbits \leftrightarrow size of twisting class
- ▶ # vertices \leftrightarrow # spanning antipodal face pairs

Twisting duality preserves:

- ▶ # antipodal face pairs

TWISTED PRISMS

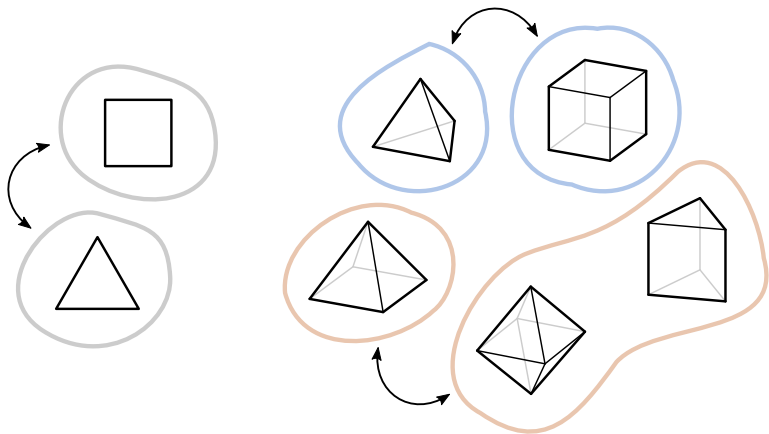


Lemma.

The twisted prism $P \boxplus P$ is 2-level if and only if P is strongly 2-level.

$P \boxplus P$	P
faces	antipodal face pairs
facets	spanning antipodal face pairs
polar duality	twisting duality
# facet types	size of twisting class

TWISTING CLASSES



Thank you.

