

BEING NON-AFRAID OF NON-CONVEXITY

– A MINIMALIST FRAMEWORK FOR GEOMETRIC PL COMPLEXES –

Martin Winter

August 7, 2026



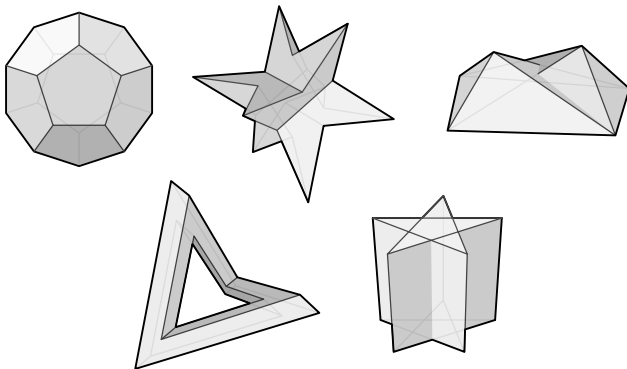
MAX PLANCK INSTITUTE
FOR MATHEMATICS IN THE SCIENCES



THE LANGUAGE OF NON-CONVEXITY

“Are your Polyhedra the same as my Polyhedra?”

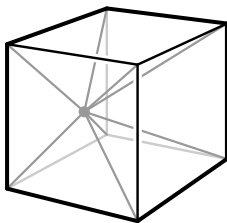
– Branko Grünbaum



MOTIVATION

- ▶ dual volumes and the canonical form
- ▶ existence of the Wachspress stress
- ▶ strong stress-flex conjecture
- ▶ spatial Maxwell-Cremona correspondence

RIGIDITY OF CONED POLYTOPE FRAMEWORKS



Theorem. (PACHYLI, PROSANOV, W., 2026)

Coned polytope frameworks are prestress stable (in particular, second-order rigid).

Ingredients:

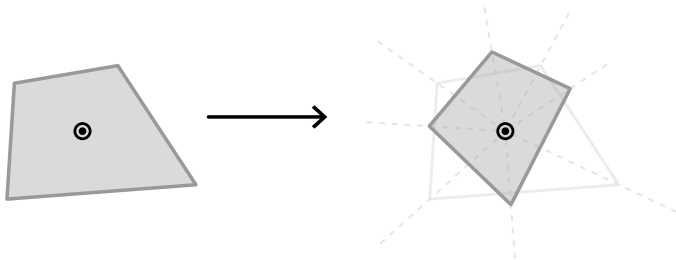
- choose a special stress $\omega : E \rightarrow \mathbb{R}$:

$$\sum_{j:i \in E} \omega_{ij}(p_j - p_i) = 0 \quad \text{for all } i \in V$$

- show that it satisfies the stress-flex conjecture

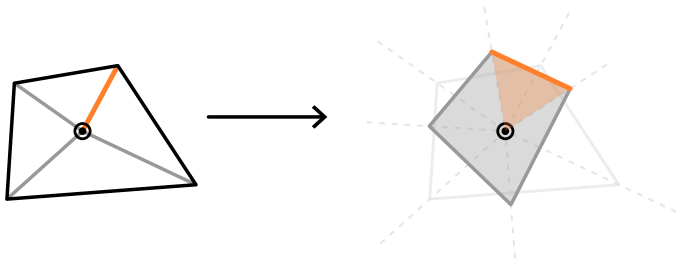
THE WACHSPRESS STRESS

$$P^\circ = \{x \in \mathbb{R}^d \mid \langle x, y \rangle \leq 1 \text{ for all } y \in P\}.$$



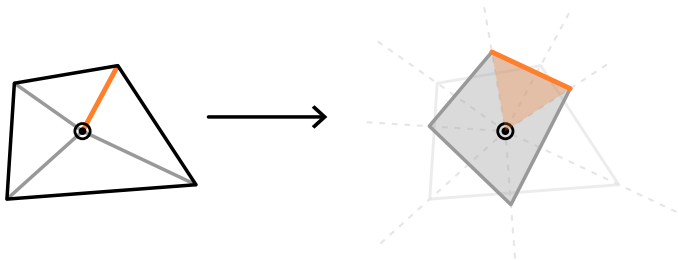
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One can show that this is a stress using **Minkowski balancing**:

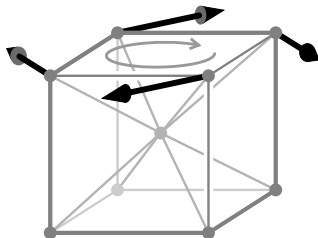
$$\sum_{\sigma \text{ codim } 1} V_\sigma n_\sigma = 0$$

STRESS FLEX CONJECTURE

Conjecture.

Given a coned polytope framework with infinitesimal flex $\dot{p}$ and stress ω , then

$$\sum_{i \neq \star} \omega_{i\star} \dot{p}_i = \dot{p}_\star.$$



GENERALIZED SCHLÄFLI FORMULAS

Given a differentiable 1-parameter family of polyhedra $P(t)$, it holds

$$\sum_e \dot{\theta}_e \operatorname{len}(e) = 0$$

GENERALIZED SCHLÄFLI FORMULAS

Given a differentiable 1-parameter family of polyhedra $P(t)$, it holds

$$\sum_{\tau \text{ codim } 2} \dot{\theta}_{\tau} \operatorname{vol}(\tau) = 0$$

GENERALIZED SCHLÄFLI FORMULAS

Given a differentiable 1-parameter family of polyhedra $P(t)$, it holds

$$\sum_{\tau \text{ codim } 2} \dot{\theta}_{\tau} \operatorname{vol}(\tau) b_{\tau} = \sum_{\sigma \text{ codim } 1} \dot{V}_{\sigma} n_{\sigma}$$

(SOUAM, SCHLENKER)

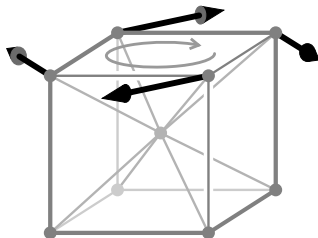
This is dual to the **stress-flex conjecture** for the **Wachspress stress**.

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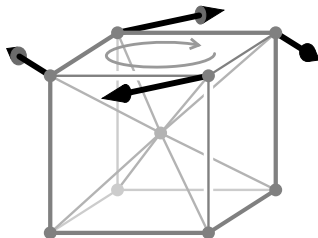


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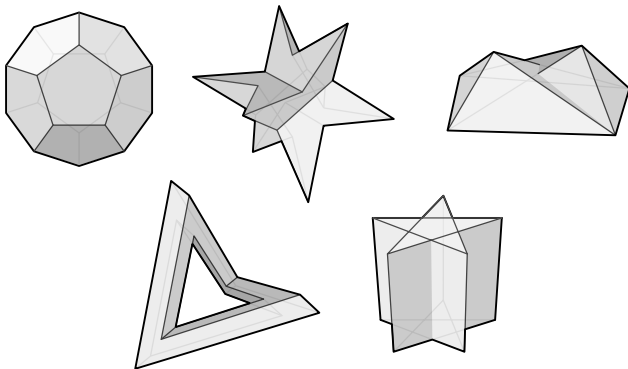
Given a coned *PL surface* framework with infinitesimal flex $\dot{p}$ and stress ω , then

$$\sum_{i \neq \star} \omega_{i\star} \dot{p}_i = \dot{p}_\star.$$



HOW MUCH NON-CONVEX DO WE WANT?

- ▶ higher genus
- ▶ self-intersecting
- ▶ faces of general combinatorics/geometry
- ▶ ...

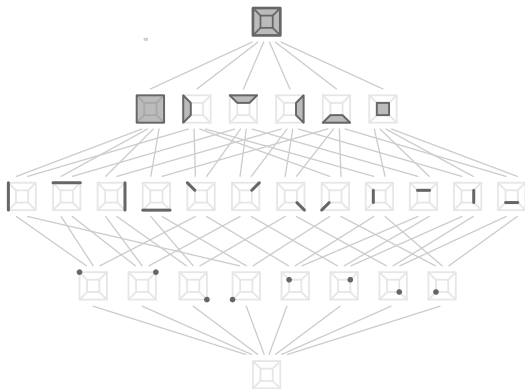
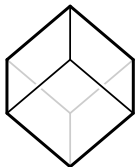


Key insight: all we need is to impose certain **coplanarities**.

THE COMBINATORIAL STRUCTURE

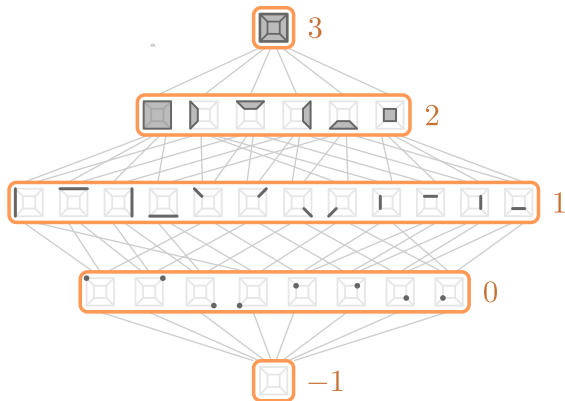
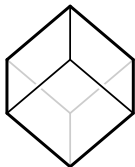
FACE LATTICES

= a **bounded graded lattice** capturing the combinatorics of a polytope.



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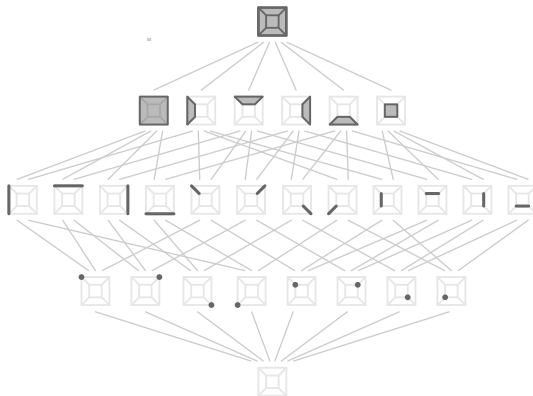
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ABSTRACT POLYTOPES (OR BETTER, ABSTRACT PL MANIFOLDS)

An **abstract polytope** is a bounded graded lattice that satisfies:

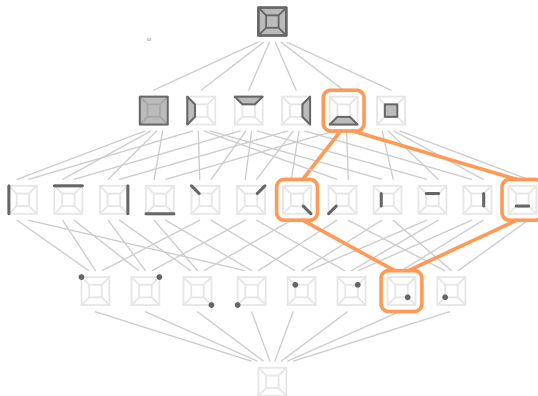
- (i) diamond property ($\simeq$ manifoldness)
- (ii) strongly flag connected (= the flag graph of every interval is connected)



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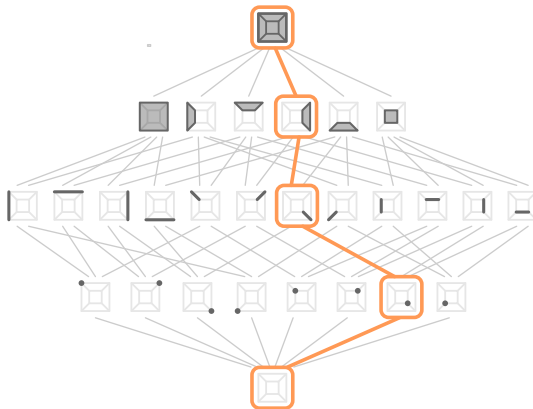
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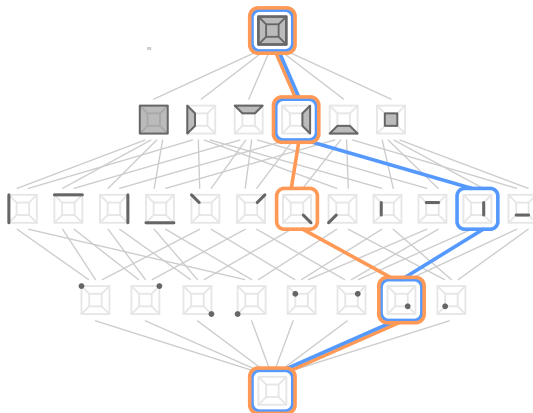
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THE GEOMETRIC REALIZATION

REALIZATIONS

An **affine realization** of $\mathcal{P}$ is a map

$$\sigma \mapsto \text{aff}(\sigma),$$

that to every face $\sigma \in \mathcal{P}$ assigns an affine subspace $\text{aff}(\sigma)$, so that

- (i) $\dim(\text{aff}(\sigma)) = \dim(\sigma)$
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That's it. There is no mention of ...

- ▶ interior/exterior,
- ▶ boundary,
- ▶ polytopes, polygons, line segments or other flat or convex pieces,
- ▶ gluing conditions,
- ▶ non-degeneracy or non-coincidence conditions

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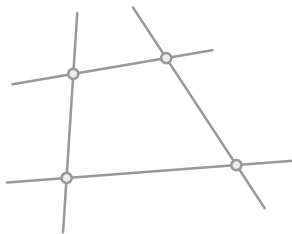
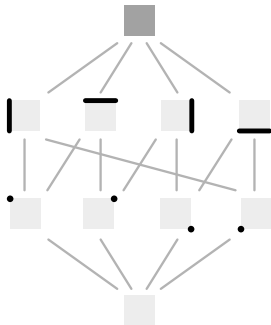
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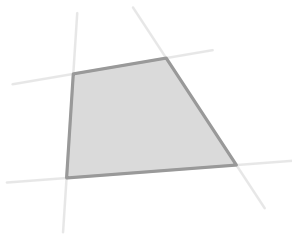
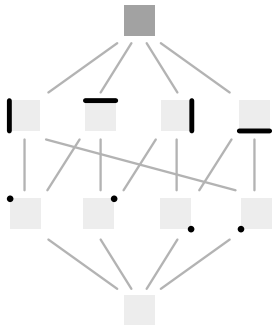
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Claim: This (almost) suffices for all we want to do.

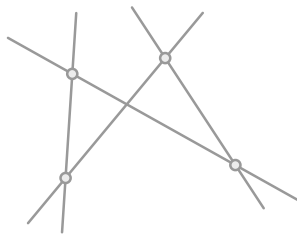
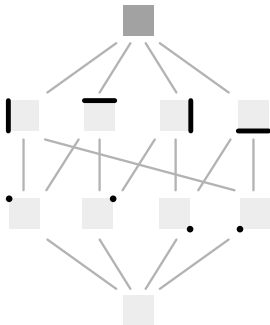
EXAMPLE



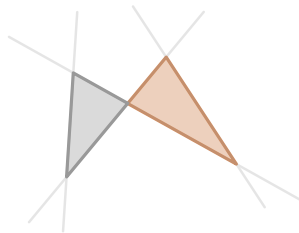
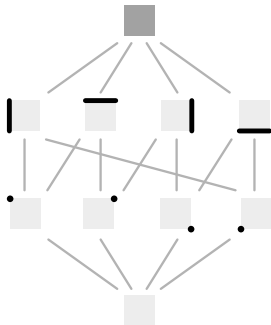
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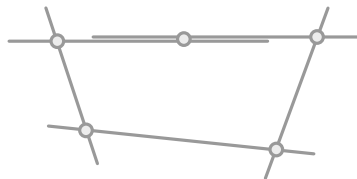
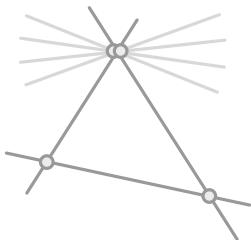
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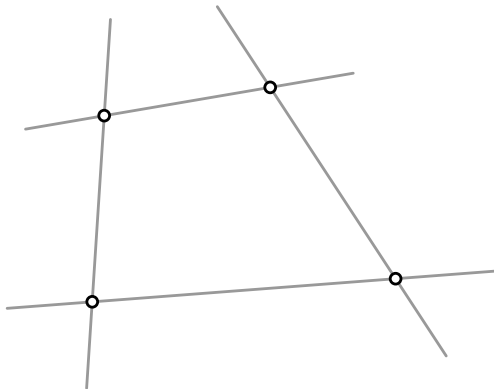


EXAMPLE (DEGENERACY IS EXPLICITLY ALLOWED)

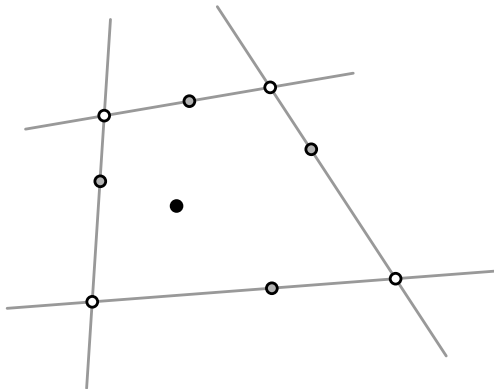


VOLUME, POLARITY AND OTHER THINGS

COMPUTING VOLUME

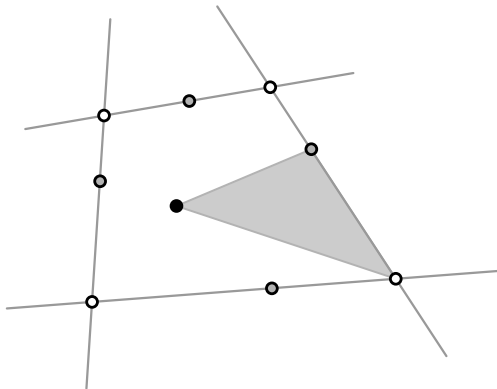


COMPUTING VOLUME



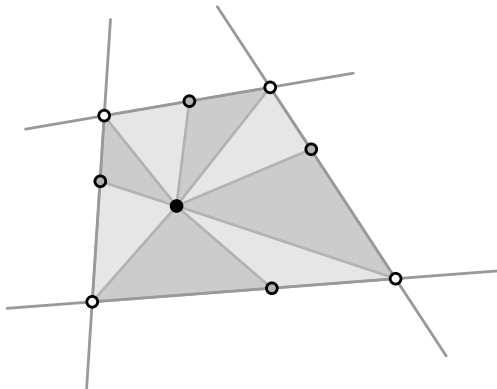
- A **pointing** maps $\sigma \mapsto x_\sigma \in \text{aff}(\sigma)$.

COMPUTING VOLUME



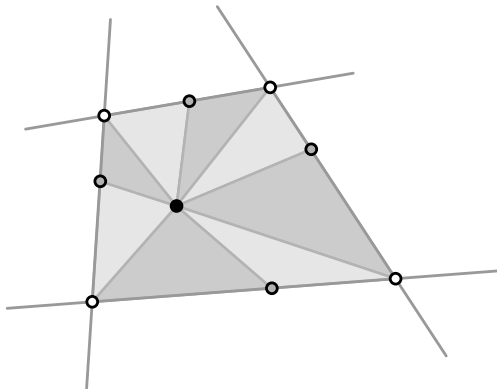
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- ▶ a flag $\sigma = \{\sigma_0 \triangleleft \sigma_1 \triangleleft \cdots \triangleleft \sigma_d\}$ yields a simplex $\Delta_\sigma := \text{conv}\{x_{\sigma_0}, \dots, x_{\sigma_d}\}$

COMPUTING VOLUME



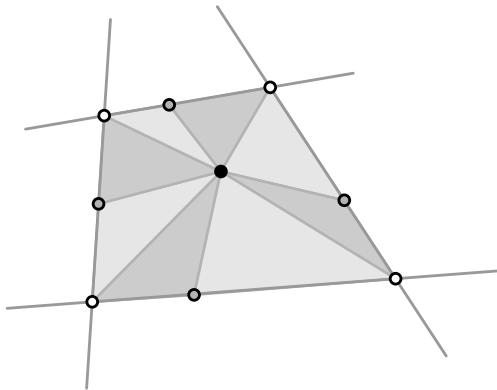
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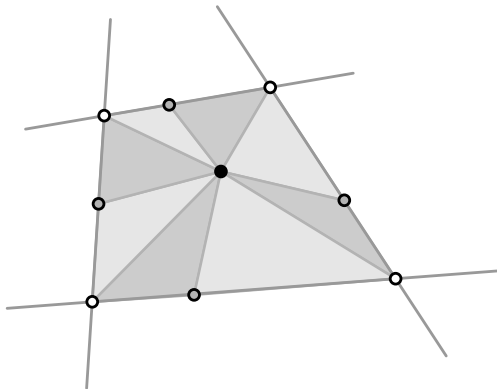
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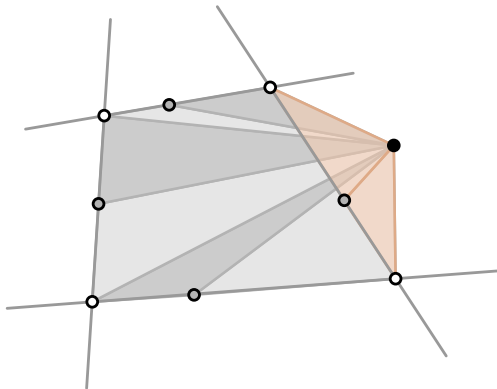
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DEFINING VOLUME



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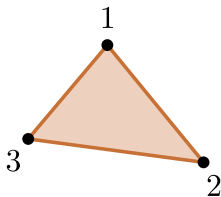
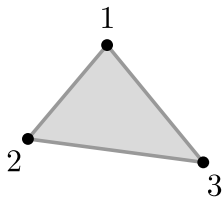
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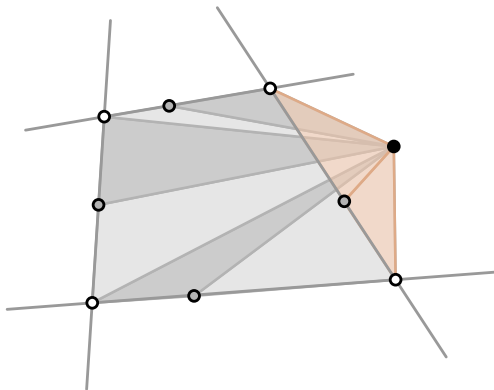
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ORIENTATION

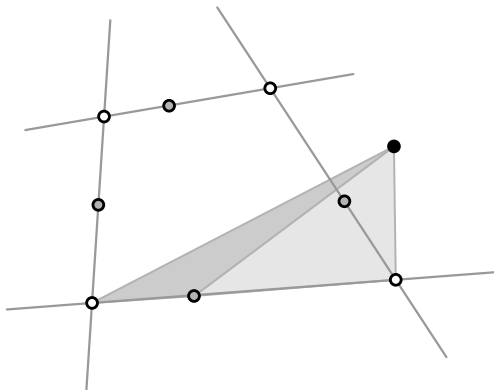
- ▶ $\mathcal{P}$ is **orientable** if the flag graph is bipartite.
- ▶ an **combinatorial orientation** of $\mathcal{P}$ is bipartition of the flag graph into classes $+1$ and -1 .
- ▶ A **geometric orientation** of $\mathbb{R}^d$ is selection of signs for d affinely independent points.



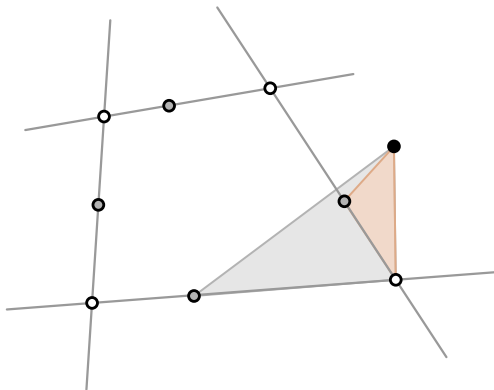
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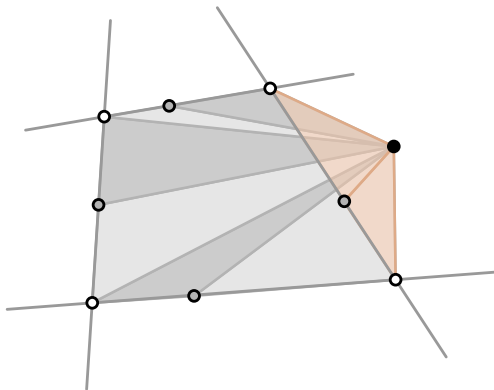
ORIENTATION



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DEFINING VOLUME WITH ORIENTATION IN MIND

$$\mathrm{vol}(\mathcal{P}, \mathrm{aff}) := \sum_{\sigma} \mathrm{sgn}(\sigma) \mathrm{sgn}(\Delta_{\sigma}) \mathrm{vol}(\Delta_{\sigma}).$$

DEFINING VOLUME WITH ORIENTATION IN MIND

A **valuation** is a map $\phi : \mathfrak{P}^d \rightarrow A$ with

$$\phi(P) + \phi(Q) = \phi(P \cap Q) + \phi(P \cup Q)$$

whenever $P, Q, P \cup Q, P \cap Q$ are convex polytope.

$$\phi(\mathcal{P}, \text{aff}) := \sum_{\sigma} \text{sgn}(\sigma) \text{sgn}(\Delta_{\sigma}) \phi(\Delta_{\sigma}) \quad .$$

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A valuation is **simple** if $\phi(P) = 0$ whenever $\dim(P) < d$. Equivalently

$$\phi(P_1 \cup \dots \cup P_n) = \phi(P_1) + \dots + \phi(P_n).$$

Theorem. (M., 2026+)

If ϕ is simple, then $\phi(\mathcal{P}, \text{aff})$ does not depend on the pointing.

MANY THINGS ARE SIMPLE VALUATION

Simple valuations:

- ▶ volume
- ▶ odd surface area measure
- ▶ Dehn invariant $\sum_e \ell_e \otimes_{\mathbb{Z}} (\theta_e/\pi)$
- ▶ dual volume $\text{vol}(P - x)^\circ$ (canonical form)
- ▶ indicator functions modulo lower dimensional polytopes
- ▶ normal vectors weighted by facet volume (on a fixed hyperplane)
- ▶ dihedral angles modulo π (on a fixed codim-2 subspace)

Valuative identities:

- ▶ Minkowski balancing
- ▶ Schläfli formula (and its extensions)
- ▶ stress-flex orthogonality

GOING BEYOND SIMPLE VALUATIONS

- ▶ Not every valuation extends to this setting
- ▶ Non-simple valuations require geometric orientations on the affine subspaces
- ▶ Some valuations do not extend uniquely
- ▶ It works for general *isometry-covariant valuations*
e.g. intrinsic volumes

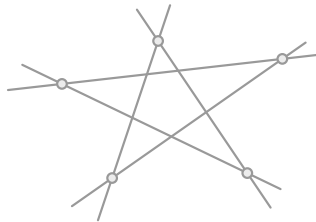
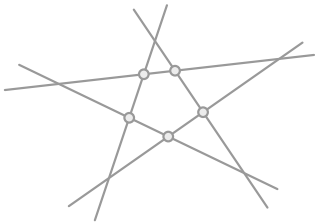
Thank you.

- ▶ Many results about PL manifolds do not rely on being convex, genus zero, being embedded, “nice” faces, ...
- ▶ Realizations of abstract polytopes via affine subspaces suffices to define a wide range of natural geometric quantities (including simple valuations) and identities.

“Second-order rigidity of coned polytope frameworks and the stress-flex conjecture from a vector-valued Schläfli formula”.

with Eleni Pachyli and Roman Prosanov. [arXiv:2607.14878](https://arxiv.org/abs/2607.14878)

EXAMPLE



PROOF OF INDEPENDENCE

Idea: Lift the valuation to a universal simple valuation, and prove it there.

- ▶ Let A be the abelian group generated by formal oriented d -simplices $[x_0, \dots, x_d]$, factored by relations:
 - ▶ $[x_{s(0)}, \dots, x_{s(d)}] = \text{sgn}(s)[x_0, \dots, x_d]$
 - ▶ $[x_0, \dots, x_d] = 0$ if the points are coplanar
 - ▶ $[\Delta] = [\Delta_1] + \dots + [\Delta_n]$ if the Δ_i tile Δ and have the same orientations.
- ▶ The map $\Delta \mapsto [\Delta]$ extends to a simple valuation $i : \mathfrak{P}^d \rightarrow A$.
- ▶ Every simple valuation factors through i : $\phi(P) = \tilde{\phi}(i(P))$.
- ▶ It suffices to prove that $i(\mathcal{P}, \text{aff})$ is independent of the pointing.
- ▶ *[technical argument about simplicial chains using the diamond property]*

POLAR DUALITY

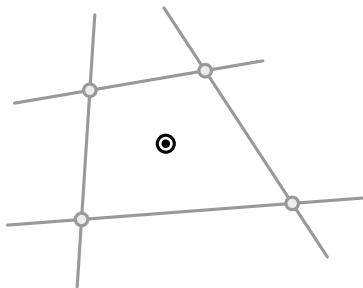
Combinatorially:

► $\dim(\sigma^\diamond) = d - \dim(\sigma) - 1$

Geometrically:

► $\text{aff}(\sigma^\diamond) := \text{aff}(\sigma)^\perp$

► $\text{dist}(\text{aff}(\sigma^\diamond)) = \text{dist}(\text{aff}(\sigma))^{-1}$



POLAR DUALITY

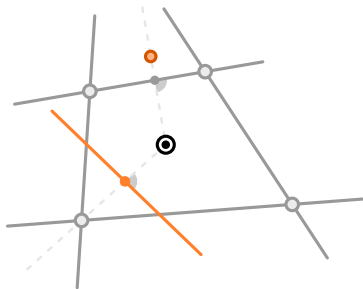
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POLAR DUALITY

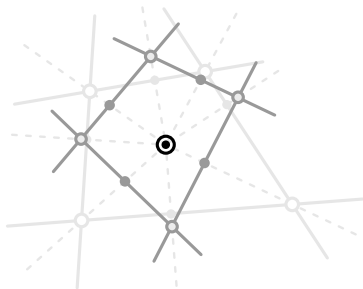
Combinatorially:

► $\dim(\sigma^\diamond) = d - \dim(\sigma) - 1$

Geometrically:

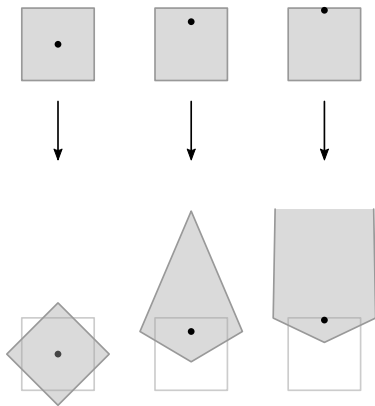
► $\text{aff}(\sigma^\diamond) := \text{aff}(\sigma)^\perp$

► $\text{dist}(\text{aff}(\sigma^\diamond)) = \text{dist}(\text{aff}(\sigma))^{-1}$



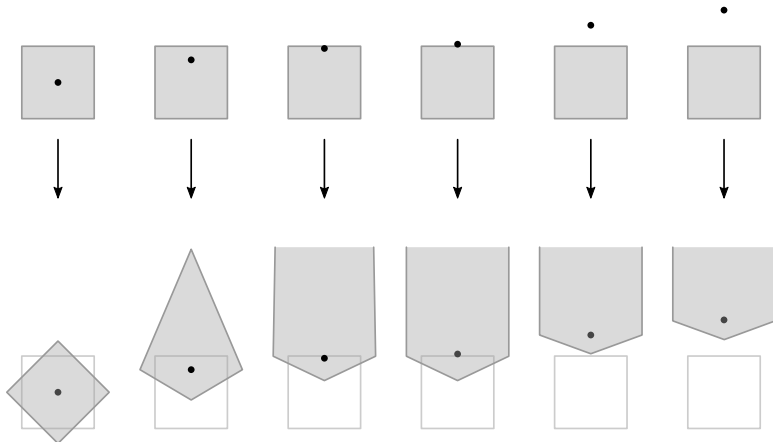
POLARITY BEYOND CONVEXITY

$$P^\circ = \{x \in \mathbb{R}^d \mid \langle x, y \rangle \leq 1 \text{ for all } y \in P\}.$$



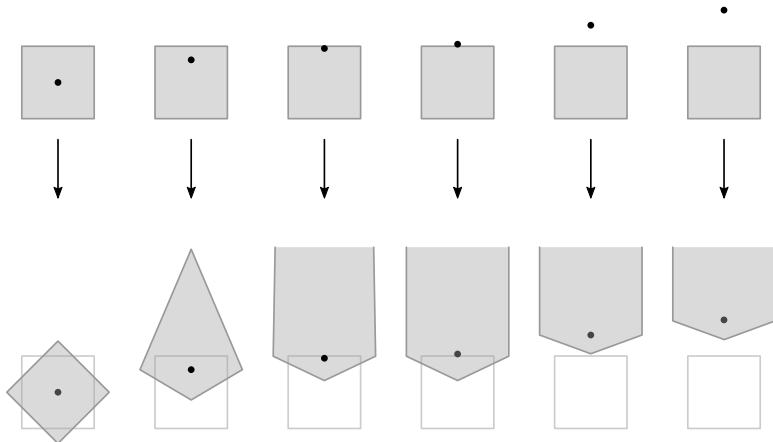
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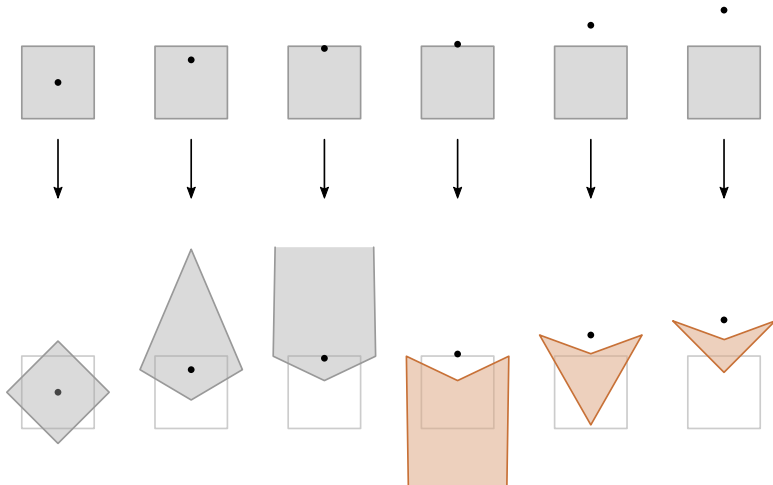
POLARITY BEYOND CONVEXITY

$$\text{vol}(P - x)^\circ = \frac{p(x)}{q(x)}$$

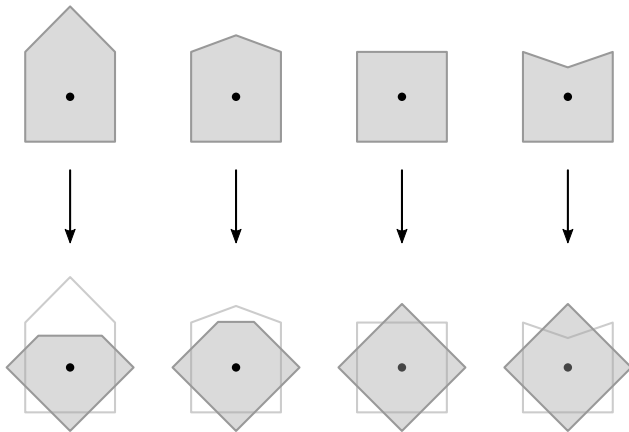


POLARITY BEYOND CONVEXITY

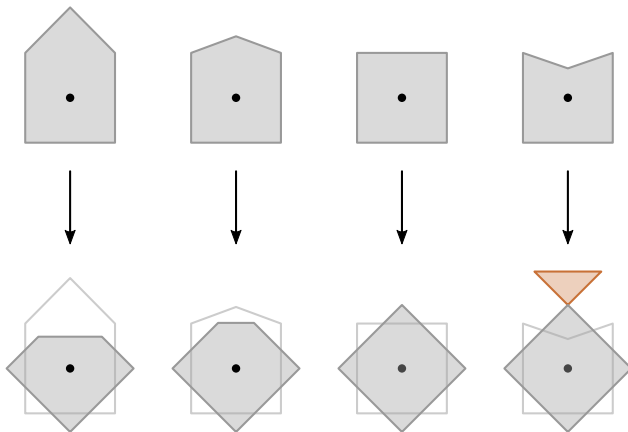
$$\text{vol}(P - x)^\circ = \frac{p(x)}{q(x)}$$



POLARITY BEYOND CONVEXITY



POLARITY BEYOND CONVEXITY



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